CENTERED LOGIC: THE ROLE OF ENTITY CENTERED SENTENCE REPRESENTATION IN NATURAL LANGUAGE INFERENCE

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ABSTRACT: We will briefly describe the role of entity centered structure (ECS) of sentences in natural language inferencing. The basic structure of sentences in discourse, generally singles out an entity, to be called center, among all those which are the arguments of the main predicate. ECS makes n-ary predicates look like monadic by temporarily masking their structure, thereby affecting the relative ease with which certain inferences are made and information is retrieved. This short paper deals with a preliminary formulation of a system designed to capture these ideas and contains several examples of how some natural language inferences can be represented in the system. Formal properties of the system are under investigation.

1. Introduction: A uniform mechanism that subsumes all inference mechanisms involved in problem solving in general may be adequate to characterize inferences in language in some sense (analogy: Turing machines characterize all computable functions); however, it will not shed much light on those mechanisms that are language relevant and presumably contribute to the efficiency of the inferencing process. In a natural language inferencing system, we are concerned with not just what inferences are made, but also how they are made (with what ease, for example). This paper is motivated by these considerations. In particular, we will be concerned with the fact that the basic structure of natural language sentences in discourse, generally, singles out an individual (entity), to be called the center, among all those which are the arguments of the main predicate. Our notion of center roughly corresponds to the linguistic notions of focus (in contrast to presupposition) and comment (in contrast to topic). We have deliberately used a new term in order to avoid precise identification with those notions and the possible resultant misunderstanding. These linguistic notions are somewhat vague and there is a great deal of confusion in the literature; further, the term "focus" is used in the AI literature in yet another sense. For the interested reader, we recommend [4], [7], and [8] for the linguistic notions, and [5], [6] for the AI notions.

The notion of center is a discourse construct; it may on occasion map on the subject of the sentence, but this need not be the case always. Such a representation can be regarded as an ascription of a property to a single individual, though the property itself may involve other individuals. For instance, in a particular context, JOHN may be the center of the sentence as in (1). JOHN HIT BILL. Underlining designates the center. It may help the reader to represent (1) in the extraposed form where the center is more clearly indicated as in (2) IT IS JOHN WHO HIT BILL. In another context, the center may be BILL, as in (3) JOHN HIT BILL. (IT IS BILL WHOM JOHN HIT.) More formally, we will represent (1) and (3) respectively as (4) and (5): (4)(JOHN x) (HIT x BILL) or (jx) (H x b); (5) (BILL y) (HIT y) or (jy) (H j y).

This entity centered structure (ECS) of natural language sentences is in sharp contrast to the structure of the usual formal language sentences which express relations among several individuals without singling out any one in particular. (Our use of the term ECS is very limited, precisely as defined here. For a wider use of this
term in the context of knowledge representation (KR), see [2] and [1].) ECS makes it easier to see the rough logical form of the sentence. n-ary predicates seem to be monadic because part of their structure is temporarily hidden (see 4 and 5 above). An individual which is not centered is essentially ignored. It can be brought into consideration only by being made a center. ECS seems to affect the relative ease with which certain inferences are made and information is retrieved. For example, given (1) IT WAS JOHN WHO HIT BILL, (jx) (H x b), it is easier to answer (2) WHO HIT BILL? (ax?) (H x b) than (3) WHOM DID JOHN HIT? (ay?) (H j y). (See also Examples 1 and 2 in Section 5.)

It is well known that inferencing is much easier in monadic predicate logic (MPL) than in full predicate logic, and that it is desirable (see, for example, [3]). It is also obvious that MPL is inadequate to handle all inferences in natural language. Nevertheless, a great deal of inferencing in natural language seems to proceed as if we are dealing with monadic predicates. It is of interest to see what mechanisms can be added to MPL which would give the added power, yet maintain the essential flavor of monadic calculus. We have been investigating some systems of this nature from the points of view of (i) their ability to capture some key properties of inference mechanisms at work in language, and (ii) their formal properties, to the extent these formal results give some insight into the structure and function of discourse constructs. In Sections 2, 3, and 4 we have presented a tentative formulation of Centered Logic (CL), and in Section 5 we have presented some sample derivations, followed by some discussion of the power of CL and some open questions.

2. Language: The formal language for centered logic differs from that of predicate logic in that the individual constants serve as variable binders in the "basic sentences," e.g., we write (jx) (H x m) for IT WAS JOHN WHO HIT MARY, and (mx) (H j x) for IT WAS MARY WHO JOHN HIT. For the purposes of this paper, a simpler notation would have sufficed. Ours was chosen with an eye towards future treatments of intensional predicates. We want to be able to distinguish between, e.g., (ax) (Bxad) (Anastasia is such that she believes Anastasia is Grand Duchess of Russia) and (ax) (Bxzd) (Anastasia is such that she believes she is Grand Duchess of Russia.)

In addition, we allow complex sentences to be built-up using truth functional connectives and quantifiers. A quantified sentence is constructed by universally or existentially generalizing on a name, i.e., by replacing all occurrences of a name by a new variable \( v \) and prefixing the sentence with \( \forall v \) or \( \exists v \). For example, we have \( (\forall y) (\forall x) \) (L x y): EVERYONE LOVES SOMEONE (OR OTHER); \( (\exists y) (\forall x) \) (L x y): EVERYONE LOVES SOMEONE (IN PARTICULAR); \( (\exists y) (\forall x) \) (L x y): SOMEONE IS LOVED BY EVERYONE. Note that the most natural reading for EVERYONE LOVES SOMEONE seems to be the one which does not require generalizing on names inside the predicate; the alternative reading is more naturally expressed by SOMEONE IS LOVED BY EVERYONE, which again does not require generalizing on predicate names. This jibes with our general view that names which are not centers are not taken as seriously as those which are.

3. Inference: We often speak of the "logical form" of a sentence as if it is unique. But we all know that the task of testing English arguments for validity can be considerably simplified if we chose the "right" representation. The trick is to expose only as much structure as is needed to see that the argument is valid: the less structure needed, the easier the inference. Predicates do have structure, but it can be provisionally masked and uncovered later only when necessary. Not all unravelings of the predicate are equally difficult. A major feature of the CL is that we ignore entities that are not centered.

Derivations are trees; each node follows from its predecessor or predecessors by one of the rules. The most important feature is that the means of introducing sentences with new centers is restricted. The centers of the premises can be thought of as the initial set of centered entities (SCE). A new entity can be introduced into SCE by making it the center of a sentence and there are only two ways in which this can be done. One is to use the change of center rule, which allows us to change a sentence into one which "says the same thing" but singles out different individual as the center. The other is the introduction of a temporary assumption with an entity that was not previously in SCE. The number of applications of change of centers and the number of individuals brought into SCE are measures of the difficulty of a derivation.

The rules are divided into four groups. 1. These rules allow the usual kinds of inferences, e.g., A\( \wedge B \) from premises A and B. Most rules in this group do not need the structure of predicates. Constraints on quantification: (a) universal instantiation will be allowed only with names for entities in SCE, i.e., we can infer (mx) (L j x) from (yy)(mx)(L y x) and an additional premise of the form (jy)(P...y...).
(6) Existential generalization will be allowed only on names for entities in SCE. Thus, we can infer \((\exists y)(x)(L x y)\) from \((x)(L x y)\) if \(y\) is SCE. Of course, we can always infer \((\exists x)(L x y)\) from \((x)(L x y)\) since the entity named \(m\) is centered. 2. These rules concern inferences which turn on the structure of the predicate, but which do not require us to recognize any names occurring in the predicate, e.g., we can derive \((\exists x)(P x) \wedge (\exists y)(Q y)\) from \((\exists x)(P x \wedge Q y)\). With these rules, we can bring out the structure of the predicate so that the sentential rules can be applied to it. 3. This group contains only one rule, the change of center rule, which allows us to infer, e.g., \((\exists x)(G x m)\) which attributes a property to John (say, the property of giving Mary (the book) "Artificial Intelligence") from another sentence \((\exists x)(G x m)\) which attributes a property to "Artificial Intelligence" (say, the property of being given by John to Mary). 4. Finally, we have rules for changing bound variables and "instantiating" centers. In the usual formal system, changing bound variables requires an application of instantiation followed by another of generalization. In our system, these rules are restricted so that the strategy might require an additional application of change of center. Since the bound variable is an accidental feature of the representation, we feel that there should be rules which allow the bound variables to be changed directly. Center instantiation is needed to get the equivalence between \((\exists x)(B x d)\) and \((\exists x)(B x d)\).

4. A Formal System (tentative version): The language of Centered Logic (CL) is defined as follows, given a formulation of first order predicate calculus (FOPC). A predicate of CL is a formula of FOPC containing at most one free variable. An atomic sentence of CL is of the form \((\exists x)(P x)\) where \(P\) is a predicate containing no free variable other than \(x\) and \(a\) is an individual constant. A sentence of CL is a member of the smallest set \(X\) containing the atomic sentences and closed under the following conditions: (i) If \(A\) and \(B\) are members of \(X\), then so are \((A \wedge B)\), \((A \vee B)\), and \((A - B)\). (ii) If \(A\) and \(B\) are members of \(X\) and \(A^2\) is a result of substituting \(x\) for all occurrences of \(a\) in \(A\), then \(\forall A\), \((\forall x)(A x)\), \((\exists x)(A x)\) are all members of \(X\). A pseudo-sentence of CL is either a sentence of CL or the result of substituting an individual variable for an individual constant in a sentence of CL. Some of the metamathematical variables we use are: \(A\), \(B\), \(C\), \(D\), ... for pseudo-sentences of CL; \(P\), \(Q\), \(R\), ... for predicates of CL; \(a\), \(b\), \(c\), ... for individual constants of FOPC; \(u\), \(v\), \(w\), \(x\), \(y\), \(z\), ... for variables of FOPC. The translations between CL and FOPC are obvious. (Note: CL could have been formulated as a kind of \(\lambda\)-calculus in which \(\lambda\) operators are used to form predicates from formulas subject to two restrictions: (i) the predicates formed are monadic; we do not have predicates of the form \(\lambda x_1 x_2 ... x_n A\). (ii) The \(\lambda\) operators cannot be nested.)

Inference Rules: The notation is that of Prawitz. Natural Deduction. Double lines indicate that the rule applies in either direction.

\[
\begin{array}{cccc}
1.1 & A & B & A \wedge B & A \wedge B & B \\
 & A \wedge B & A & A & B \\
1.2 & A & B & A \vee B & A \vee B & (A) (B) \\
 & A \vee B & A & B & (A) (B) \\
1.3 & (A) & A & A \wedge B & A \wedge B & (B) \\
 & B & A \wedge B & A \wedge B & (B) \\
1.4 & (A) (B) & (A) (B) & (A) (B) & (A) (B) & (A) & (B) \\
 & (A) & (B) & (A) & (B) & (A) & (B) \\
1.5 & A & (A x) (P X) & (A x) (P X) & (A x) (P X) & (A x) (P X) & (A x) (P X) \\
 & (A x) (P X) & (A x) (P X) & (A x) (P X) & (A x) (P X) & (A x) (P X) & (A x) (P X) \\
1.6 & (A x) & (A x) (P X) & (A x) (P X) & (A x) (P X) & (A x) (P X) & (A x) (P X) \\
 & (A x) & (A x) (P X) & (A x) (P X) & (A x) (P X) & (A x) (P X) & (A x) (P X) \\
\end{array}
\]

Restrictions: In the first rule in (1.5), \(x\) must occur in any assumption on which \(A\) depends. In the second (1.6) rule, a \(x\) must not occur in \((\forall x)(A)\), \(B\), or any assumption on which the upper occurrence of \(B\) depends except \(A^x\).

2. Predicate Decomposition

\[
\begin{array}{cccc}
2.1 & (A x) (P \wedge Q) & (A x) (P \wedge Q) & (A x) (P \wedge Q) & (A x) (P \wedge Q) \\
 & (A x) (P \wedge Q) & (A x) (P \wedge Q) & (A x) (P \wedge Q) & (A x) (P \wedge Q) \\
2.2 & (A x) (P \vee Q) & (A x) (P \vee Q) & (A x) (P \vee Q) & (A x) (P \vee Q) \\
 & (A x) (P \vee Q) & (A x) (P \vee Q) & (A x) (P \vee Q) & (A x) (P \vee Q) \\
2.3 & (A x) (P \rightarrow Q) & (A x) (P \rightarrow Q) & (A x) (P \rightarrow Q) & (A x) (P \rightarrow Q) \\
 & (A x) (P \rightarrow Q) & (A x) (P \rightarrow Q) & (A x) (P \rightarrow Q) & (A x) (P \rightarrow Q) \\
2.4 & (A x) (\neg P) & (A x) (\neg P) & (A x) (\neg P) & (A x) (\neg P) \\
 & (A x) (\neg P) & (A x) (\neg P) & (A x) (\neg P) & (A x) (\neg P) \\
2.5 & (A x) (\exists y) (P (x)) & (A x) (\exists y) (P (x)) & (A x) (\exists y) (P (x)) & (A x) (\exists y) (P (x)) \\
 & (A x) (\exists y) (P (x)) & (A x) (\exists y) (P (x)) & (A x) (\exists y) (P (x)) & (A x) (\exists y) (P (x)) \\
\end{array}
\]

Provided \(x \neq y\).
3. Change of Topics

2.1 \[
\frac{(ax)(P)}{(bx)(\langle a \rangle z)}
\]

4. Bound Variables

4.1 \[
\frac{(y)(A)}{(y)(A')}
\]

4.2 \[
\frac{(z)(A)}{(z')(A')}
\]

4.3 \[
\frac{(ax)(A)}{(ax)(A')}
\]

4.4 \[
\frac{(ax)(A)}{(ax)(A')}
\]

where \(A'\) is the result of
replacing one or more free
occurrences of \(x\) in \(A\) by \(a\).

5. Sample Derivations:

(1a) Everyone at the reception was thanked (T)
by John (J). Mary (M) was at the reception (R).
(therefore) Mary was thanked by John.

\[
\begin{align*}
(y)(x)(Rx \rightarrow Jx) & \quad (m)(x)(Rx) \\
(m')(x)(Rx & \rightarrow Tjx) \\
(m')(x)(Rx) & \rightarrow (m')(x)(Tjx) \\
(m')(x)(Rx) & \\
\end{align*}
\]

(1b) John thanked everyone at the reception.
Mary was at the reception. (therefore) John
thanked Mary.

\[
\begin{align*}
(j)(x)(y)(Ry \rightarrow Tx) & \quad (m)(x)(Rx) \\
(y)(j)(x)(Ry \rightarrow Txy) & \quad (m')(x)(Rx) \\
(j)(x)(Rm \rightarrow Txm) & \quad (m)(x)(Rx \rightarrow Tjx) \\
(m)(x)(Rx) & \rightarrow (m')(x)(Tjx) \\
(m')(x)(Rx) & \\
\end{align*}
\]

(1b) is a more difficult inference than (1a)
because it requires two applications of the
change of center rule (marked by double lines).

Neither the inference in (1a) nor in (1b)
requires bringing into SCE (Set of Centered
Entities) a new entity which was not a center of
a premise. This contrasts with the following
derivation (2) which requires two new entities to
be brought into SCE. The inherent difficulty of
the inference in (2) below is due to this require-
ment, at least with respect to the CL.

(2) There is a house in which everyone lives
(\(L\) (therefore) Everyone lives in houses.

\[
\begin{align*}
(y)(z)(Lxy) & \quad (b)(x)(P) \\
(y)(z)(Lxy) & \quad (b)(x)(P \rightarrow Lby) \\
(y)(z)(Lxy) & \quad (b)(x)(Lx)(ay)(Lzy) \\
(y)(z)(Lxy) & \quad (b)(x)(Lx)(ay)(Lzy) \\
(y)(z)(Lxy) & \quad (b)(x)(Lx)(ay)(Lzy) \\
(y)(z)(Lxy) & \quad (b)(x)(Lx)(ay)(Lzy) \\
\end{align*}
\]

* Temporary assumptions discharged later.

It is not difficult to prove that the rules
given for CL are complete in the sense that \(A\)
can be derived from \(\Gamma\) whenever the * translation
(from CL to FOPC) of \(F\) is classically derivable
from the * translation of \(\Gamma\). To see this, note
that (i) the \(\circ\)-translation (from FOPC to CL) of
the classical natural deduction rules are all
derivable from our rules, and (ii) \(A^*\) is
provable equivalent to \(A\). However, we need both
the change of center and the introduction of
temporary assumptions with new centers to get
this result. (The former is needed to establish
the equivalence of \(A\) and \(A^*\) and the latter to
show that translations of the classical quanti-
plier rules are derivable.) If premises and con-
clusions have the same center, then no change of
center rules are needed. For, if \(a\) is the center
of premises and conclusions, then by translating
each classical \(\circ\) as \((ax)\langle a \rangle\) and using the predi-
cate decomposition rules, the classical derivation
of the *-translation of the argument can be
converted to a CL derivation without center
changes. Hence, the number of center changes
need never exceed the number of premises whose
centers are distinct from the centers of the
conclusion. (The role of center change becomes
more interesting, and more difficult to under-
stand, when quantification into the predicate is
prohibited. As we mentioned previously, there
is some reason to believe that this restriction
is appropriate.)

Although a restriction could be placed on the
number of center changes, no such limit can be placed on the number of new entities which must be brought into SCE. The premise in Example (2) contains no individual constant at all. It is clear from the rules that from such formulas alone we can only derive equivalent formulas and tautologous consequences. Hence, the introduction of new centers in Example (2) was unavoidable.

Thus, the logic resulting from a restriction on the number of centered entities is less powerful than one without, in particular, the logic with zero centered entities is less powerful than the logic with one centered individual. General questions about the number of new centers required for an inference and the decidability of the class of inferences with small number of centers remain open.

References (a few selected ones):


